

# A variational derivation of a class of BFGS-like methods

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## ABSTRACT

We provide a maximum entropy derivation of a new family of BFGS-like methods. Similar results are then derived for block BFGS methods. This also yields an independent proof of a result of Fletcher 1991 and its generalisation to the block case.

## KEYWORDS

Quasi-Newton method, BFGS method, maximum entropy problem, block BFGS.

## 1. Introduction

Suppose  $f : \mathbb{R}^n \rightarrow \mathbb{R}$  is a  $C^2$  function to be minimized. Then Newton’s iteration is

$$x_{k+1} = x_k - [H(x_k)]^{-1} \nabla f(x_k), \quad k \in \mathcal{N}, \quad (1)$$

where  $H(x_k) = \nabla^2 f(x_k)$  is the Hessian of  $f$  at the point  $x_k$ . In quasi-Newton methods, one employs instead an approximation  $B_k$  of  $H(x_k)$  to avoid the costly operations of computing, storing and inverting the Hessian ( $B_0$  is often taken to be the identity  $I_n$ ). These methods appear to perform well even in nonsmooth optimization, see [1]. Instead of (1), one uses

$$x_{k+1} = x_k - \alpha_k B_k^{-1} \nabla f(x_k), \quad \alpha_k > 0, \quad k \in \mathcal{N}, \quad (2)$$

with  $\alpha_k$  chosen by a line search, imposing the *secant* equation

$$y_k = B_{k+1} s_k, \quad (3)$$

where

$$y_k := \nabla f(x_k + s_k) - \nabla f(x_k), \quad s_k := \Delta x_k = x_{k+1} - x_k.$$

The secant condition is motivated by the expansion

$$\nabla f(x_k + s_k) \approx \nabla f(x_k) + H(x_k)s_k. \quad (4)$$

For  $n > 1$ ,  $B_{k+1}$  satisfying (3) is underdetermined. Various methods are used to find a symmetric  $B_{k+1}$  that satisfies the secant equation (3) and is closest in some metric to the current approximation  $B_k$ . In several methods,  $B_{k+1}$  or its inverse is a rank one or two update of the previous estimate [2].

Since for a strongly convex function the Hessian  $H(x_k)$  is a symmetric positive definite matrix, we can think of its approximation  $B_k$  as a covariance of a zero-mean, multivariate Gaussian distribution. Recall that in the case of two zero-mean multivariate normal distributions  $p, q$  with nonsingular  $n \times n$  covariance matrixes  $P, Q$ , respectively, the relative entropy (divergence, Kullback-Leibler index) can be derived in closed form

$$\mathbb{D}(p||q) = \int \log \frac{p(x)}{q(x)} p(x) dx = \frac{1}{2} [\log \det (P^{-1}Q) + \text{tr}(Q^{-1}P) - n].$$

Since  $P^{-1}$  and  $Q^{-1}$  are the natural parameters of the Gaussian distributions, we write

$$\mathbb{D}(P^{-1}||Q^{-1}) = \frac{1}{2} [\log \det (P^{-1}Q) + \text{trace}(Q^{-1}P) - n] \quad (5)$$

## 2. A maximum entropy problem

Consider minimizing  $\mathbb{D}(B^{-1}||B_k^{-1})$  over symmetric, positive definite  $B$  subject to the secant equation

$$B^{-1}y_k = s_k. \quad (6)$$

In [3], Fletcher indeed showed that the solution to this variational problem is provided by the BFGS iterate thereby providing a variational characterization for it alternative to Goldfarb's classical one [4], [2, Section 6.1]. We take a different approach leading to a family of BFGS-like methods.

First of all, observe that  $B^{-1}y_k$  must be the given vector  $s_k$ . Thus, it seems reasonable that  $B_{k+1}^{-1}$  should approximate  $B_k^{-1}$  only in directions different from  $y_k$ . We are then led to consider the following new problem

$$\min_{\{B=B^T, B>0\}} \mathbb{D}(B^{-1}||P_k^T B_k^{-1} P_k) \quad (7)$$

subject to (6), where  $P_k$  is a rank  $n - 1$  matrix satisfying  $P_k y_k = 0$ , subject to the secant equation (6). One possible choice for  $P_k$  is the orthogonal projection

$$P_k = I_n - \frac{y_k y_k^T}{y_k^T y_k} = I_n - \Pi_{y_k}.$$

Since  $P_k B_k^{-1} P_k$  is singular, however, (7) does not make sense. Thus, to regularize the problem, we replace  $P_k$  with the nonsingular, positive definite matrix  $P_k^\epsilon = P_k + \epsilon I_n$ .

The Lagrangian for this problem is

$$\begin{aligned} \mathcal{L}(B, \lambda) &= \frac{1}{2} [\log \det (B^{-1} (P_k^\epsilon)^{-1} B_k P_k^\epsilon) + \text{tr} (P_k^\epsilon B_k^{-1} P_k^\epsilon B) - n] + \lambda_k^T [B s_k - y_k] = \\ &= \frac{1}{2} \left[ \log \det (B^{-1} B_k) + \frac{1}{2} \log \det ((P_k^\epsilon)^{-2}) + \text{tr} (P_k^\epsilon B_k^{-1} P_k^\epsilon B) - n \right] + \lambda_k^T [B s_k - y_k]. \end{aligned}$$

Observe that the term

$$\frac{1}{2} \log \det ((P_k^\epsilon)^{-2})$$

does not depend on  $B$  and therefore plays no role in the variational analysis. To compute the first variation of  $\mathcal{L}$  in direction  $\delta B$ , we first recall a simple result. Consider the map  $J$  defined on nonsingular,  $n \times n$  matrices  $M$  by  $J(M) = \log |\det[M]|$ . Let  $\delta J(M; \delta M)$  denote the directional derivative of  $J$  in direction  $\delta M \in \mathbb{R}^{n \times n}$ . We then have the following result :

**Lemma 2.1.** [5, Lemma 2] *If  $M$  is nonsingular then, for any  $\delta M \in \mathbb{R}^{n \times n}$ ,*

$$\delta J(M; \delta M) = \text{trace}[M^{-1} \delta M].$$

Observe also that any positive definite matrix  $B$  is an interior point in the cone  $\mathcal{C}$  of positive semidefinite matrices in any symmetric direction  $\delta B \in \mathbb{R}^{n \times n}$ . Imposing  $\delta \mathcal{L}(B, \lambda; \delta B) = 0$  for all such  $\delta B$ , we get, in view of Lemma 2.1,

$$\text{trace} [(-(B_{k+1}^\epsilon)^{-1} + P_k^\epsilon B_k^{-1} P_k^\epsilon + 2s_k \lambda_k^T) \delta B] = 0, \quad \forall \delta B,$$

which gives

$$(B_{k+1}^\epsilon)^{-1} = P_k^\epsilon B_k^{-1} P_k^\epsilon + 2s_k \lambda_k^T. \quad (8)$$

As  $\epsilon \searrow 0$ , we get the iteration

$$B_{k+1}^{-1} = P_k B_k^{-1} P_k + 2s_k \lambda_k^T. \quad (9)$$

Since  $P_k y_k = 0$ , in order to satisfy the secant equation

$$B_{k+1}^{-1} y_k = s_k.$$

it suffices to choose the multiplier  $\lambda_k$  so that

$$2\lambda_k^T y_k = 1.$$

We need, however, to also guarantee symmetry and positive definiteness of the solution. We are then led to choose  $\lambda_k$  as

$$\lambda_k = \frac{s_k}{2y_k^T s_k}. \quad (10)$$

Finally, notice that, under the *curvature* assumption

$$y_k^T s_k > 0, \quad (11)$$

if  $B_k > 0$ , indeed  $B_{k+1}$  in (9) is symmetric, positive definite justifying the previous calculations. We have therefore established the following result.

**Theorem 2.2.** *Assume  $B_k > 0$  and  $y_k^T s_k > 0$ . A solution  $B^*$  of*

$$\min_{\{B=B^T, B>0\}} \mathbb{D}(B^{-1} || P_k^T B_k^{-1} P_k),$$

*subject to constraint (6), in the regularized sense described above, is given by*

$$(B^*)^{-1} = \left( I_n - \frac{y_k y_k^T}{y_k^T y_k} \right) B_k^{-1} \left( I_n - \frac{y_k y_k^T}{y_k^T y_k} \right) + \frac{s_k s_k^T}{y_k^T s_k}. \quad (12)$$

### 3. BFGS-like methods

From Theorem 2.2, we get the following quasi-Newton iteration:

$$x_{k+1} = x_k - \alpha_k B_k^{-1} \nabla f(x_k), \quad x_0 = \bar{x}, \quad (13)$$

$$B_{k+1}^{-1} = \left( I_n - \frac{y_k y_k^T}{y_k^T y_k} \right) B_k^{-1} \left( I_n - \frac{y_k y_k^T}{y_k^T y_k} \right) + \frac{s_k s_k^T}{y_k^T s_k}, \quad B_0 = I_n. \quad (14)$$

Note that, for limited-memory iterations, this method has the same storage requirement as standard limited-memory BFGS, say  $(s_j, y_j)$ ,  $j = k, k-1, \dots, k-m+1$ . Now let  $v_k \in \mathbb{R}^n$  be any vector not orthogonal to  $y_k$ . Then

$$P_k(v_k) := \frac{y_k v_k^T}{y_k^T v_k} \quad (15)$$

is an oblique projection onto  $y_k$ . Employing  $P_k(v_k)$  and its transpose in place of  $\Pi_{y_k}$  in (7) and performing the variational analysis after regularisation, we get a BFGS-like iteration

$$B_{k+1}^{-1} = (I_n - P_k(v_k))^T B_k^{-1} (I_n - P_k(v_k)) + \frac{s_k s_k^T}{y_k^T s_k} \quad (16)$$

In particular, if  $v_k = s_k$ , the corresponding oblique projection is

$$P_k(s_k) = \frac{y_k s_k^T}{y_k^T s_k}.$$

In such case, (16) is just the standard (BFGS) iteration for the inverse approximate Hessian

$$B_{k+1}^{-1} = \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} \right)^T B_k^{-1} \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} \right) + \frac{s_k s_k^T}{y_k^T s_k}. \quad (17)$$

Here  $T_k = I_n - P_k(s_k)$  is a rank  $n - 1$  matrix satisfying  $T_k y_k = 0$  as is  $I - \Pi_{y_k}$ . We now get an alternative derivation of Fletcher's result [3].

**Corollary 3.1.** *Assume  $B_k > 0$  and  $y_k^T s_k > 0$ . A solution  $B^*$  of*

$$\min_{\{B=B^T, B>0\}} \mathbb{D}(B^{-1} \| B_k^{-1}),$$

*subject to constraint (6) is given by the standard (BFGS) iteration (17).*

**Proof.** We show that in the limit, as  $\epsilon \searrow 0$ ,  $\mathbb{D}(B^{-1} \| B_k^{-1})$  and  $\mathbb{D}\left(B^{-1} \left\| \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} + \epsilon I_n \right)^T B_k^{-1} \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} + \epsilon I_n \right) \right\| \right)$  only differ by terms not depending on  $B$ . Indeed,

$$\begin{aligned} & \mathbb{D}\left(B^{-1} \left\| \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} + \epsilon I_n \right)^T B_k^{-1} \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} + \epsilon I_n \right) \right\| \right) \\ &= \frac{1}{2} \left\{ \log \det(B^{-1} B_k) + \log \det \left[ \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} + \epsilon I_n \right)^{-1} \left( I_n - \frac{y_k s_k^T}{y_k^T s_k} + \epsilon I_n \right)^{-T} \right] \right. \\ & \quad \left. + \text{trace} \left[ \left( (1 + \epsilon) I_n - \frac{y_k s_k^T}{y_k^T s_k} \right)^T B_k^{-1} \left( (1 + \epsilon) I_n - \frac{y_k s_k^T}{y_k^T s_k} \right) B \right] - n \right\} \end{aligned}$$

Note that, by the circulant property of the trace,

$$\text{trace} \left[ -\frac{s_k y_k^T}{y_k^T s_k} B_k^{-1} (1 + \epsilon) B \right] = \text{trace} \left[ -B \frac{s_k y_k^T}{y_k^T s_k} B_k^{-1} (1 + \epsilon) \right]$$

It now suffices to observe that, for symmetric matrices  $B$  satisfying (6)  $B s_k = y_k$ , the products

$$B \frac{s_k y_k^T}{y_k^T s_k} = \frac{y_k s_k^T}{y_k^T s_k} B = \frac{y_k y_k^T}{y_k^T s_k}$$

are independent of  $B$ . □

Iterations (13)-(14) and (13)-(16) are expected to enjoy the same convergence properties as the canonical BFGS method [2, Chapter 6]. They can, in principle, be applied also to nonsmooth cases along the lines of [1] with an exact line search to compute  $\alpha_k$  at each step.

#### 4. Block BFGS-like methods

In some large dimensional problems, it is prohibitive to calculate the full gradient at each iteration. Consider for instance *deep neural networks*. A deep network consists of a nested composition of a linear transformation and a nonlinear one  $\sigma$ . In the learning phase of a deep network, one compares the predictions  $y(x, \xi^i)$  for the input sample  $\xi^i$  with the actual output  $y^i$ . This is done through a cost function  $f_i(x)$ , e.g.

$$f_i(x) = \|y^i - y(x; \xi^i)\|^2.$$

The goal is to learn the *weights*  $x$  through minimization of the empirical loss function

$$f(x) = \frac{1}{N} \sum_{i=1}^N f_i(x).$$

In modern datasets,  $N$  can be in the millions and therefore calculation of the full gradient  $\frac{1}{N} \sum_{i=1}^N \nabla f_i(x)$  at each iteration to perform gradient descent is unfeasible. One can then resort to *stochastic gradients* by sampling uniformly from the set  $\{1, \dots, N\}$  the index  $i_k$  where to compute the gradient at iteration  $k$ . In alternative, one can also average the gradient over a set of randomly chosen samples called a “mini-batch”. In [6], a so-called block BFGS was proposed. Let  $S_k$  be a *sketching matrix* of directions [6] and let  $\mathcal{T} \subset [N]$ . Rather than taking differences of random gradients, one computes the action of the sub-sampled Hessian on  $S_k$  as

$$Y_k := \frac{1}{|\mathcal{T}|} \sum_{i \in \mathcal{T}} \nabla^2 f_i(x_k) S_k$$

To update  $B_k^{-1}$ , we can now consider the problem

$$\min_{\{B=B^T, B>0\}} \mathbb{D} (B^{-1} \| P_k^T B^{-1} P_k) \quad (18)$$

where  $I - P_k$  projects onto the space spanned by the columns of  $Y_k$ , subject to the block-secant equation

$$B^{-1} Y_k = S_k. \quad (19)$$

Again, one possible choice for  $S_k$  is  $I - \Pi_{Y_k}$  where  $\Pi_{Y_k} = Y_k (Y_k^T Y_k)^{-1} Y_k^T$  is the orthogonal projection. The same variational argument as in Section 2 leads to the iteration

$$B_{k+1}^{-1} = (I - \Pi_{Y_k}) B_k^{-1} (I - \Pi_{Y_k}) + S_k (S_k^T Y_k)^{-1} S_k^T. \quad (20)$$

Another choice for  $P_k$  is the oblique projection  $I - Y_k (S_k^T Y_k)^{-1} S_k^T$  leading to the iteration in [6]

$$B_{k+1}^{-1} = (I - Y_k (S_k^T Y_k)^{-1} S_k^T)^T B_k^{-1} (I - Y_k (S_k^T Y_k)^{-1} S_k^T) + S_k (S_k^T Y_k)^{-1} S_k^T. \quad (21)$$

We then obtain a variational characterisation of the iteration (21) alternative to the one of [6, Appendix A] and generalizing Fletcher [3].

**Corollary 4.1.** *Assume  $B_k > 0$  and  $S_k^T Y_k > 0$ . A solution  $B^*$  of*

$$\min_{\{B=B^T, B>0\}} \mathbb{D}(B^{-1} || B_k^{-1}),$$

*subject to constraint (19) is given by  $B_{k+1}$  in (21).*

The proof is analogous to the proof of Corollary 3.1.

## 5. Numerical Experiments

The algorithm (13)-(14) has the form:

```

1: procedure BFGS-LIKE( $f, Gf, x_0, tolerance$ )
2:    $B \leftarrow I_d$   $\triangleright d$  is the dimension of  $x_0$  and  $I_d$  is the identity in  $R^d$ 
3:    $x \leftarrow x_0$ 
4:   for  $n = 1, \dots, MaxIterations$  do
5:      $y \leftarrow Gf(x)$ 
6:     if  $\|y\| < tolerance$  then
7:       break
8:      $SearchDirection \leftarrow -By$ 
9:      $\alpha \leftarrow LineSearch(f, GF, x, SearchDirection)$ 
10:     $\Delta x \leftarrow \alpha SearchDirection$ 
11:     $S \leftarrow I_d - \frac{yy^T}{y^T y}$ 
12:     $B \leftarrow S^T B S + \frac{\Delta x \Delta x^T}{y^T \Delta x}$ 
13:     $x \leftarrow x + \Delta x$ 
14: return  $x$ 

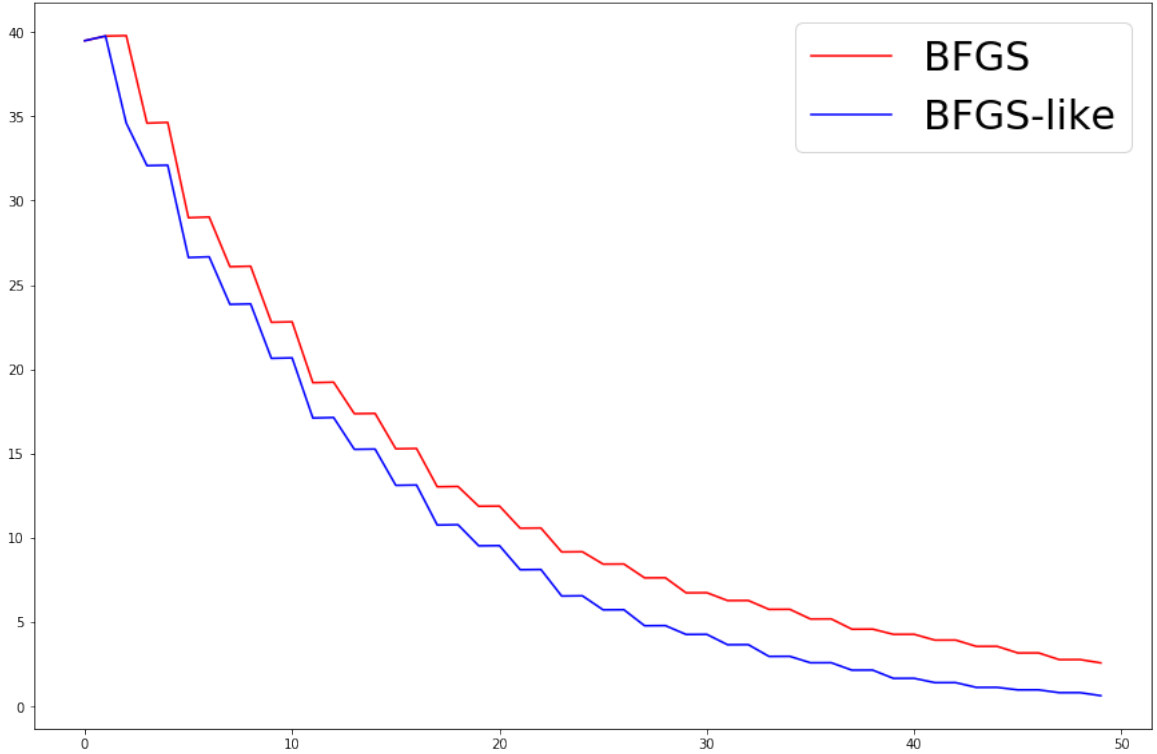
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Algorithm 1: BFGS-like algorithm (13)-(14)

While the effectiveness of the BFGS-like algorithms introduced in Section 3 needs to be tested on a significant number of large scale benchmark problems, we provide below two examples where the BFGS-like algorithm (13)-(14) appears to perform better than standard BFGS. Consider the strictly convex function  $f$  on  $\mathbb{R}^2$

$$f(x_1, x_2) = e^{x_1-1} + e^{-x_2+1} + (x_1 - x_2)^2$$

whose minimum point is  $x^* \approx (0.8, 1.2)$ . Take as starting point:  $(5, -7)$ . Figure 1 illustrates the decay of the error  $\|x^n - x^*\|_2$  over 50 iterations for the classical BFGS and for algorithm (13)-(14).



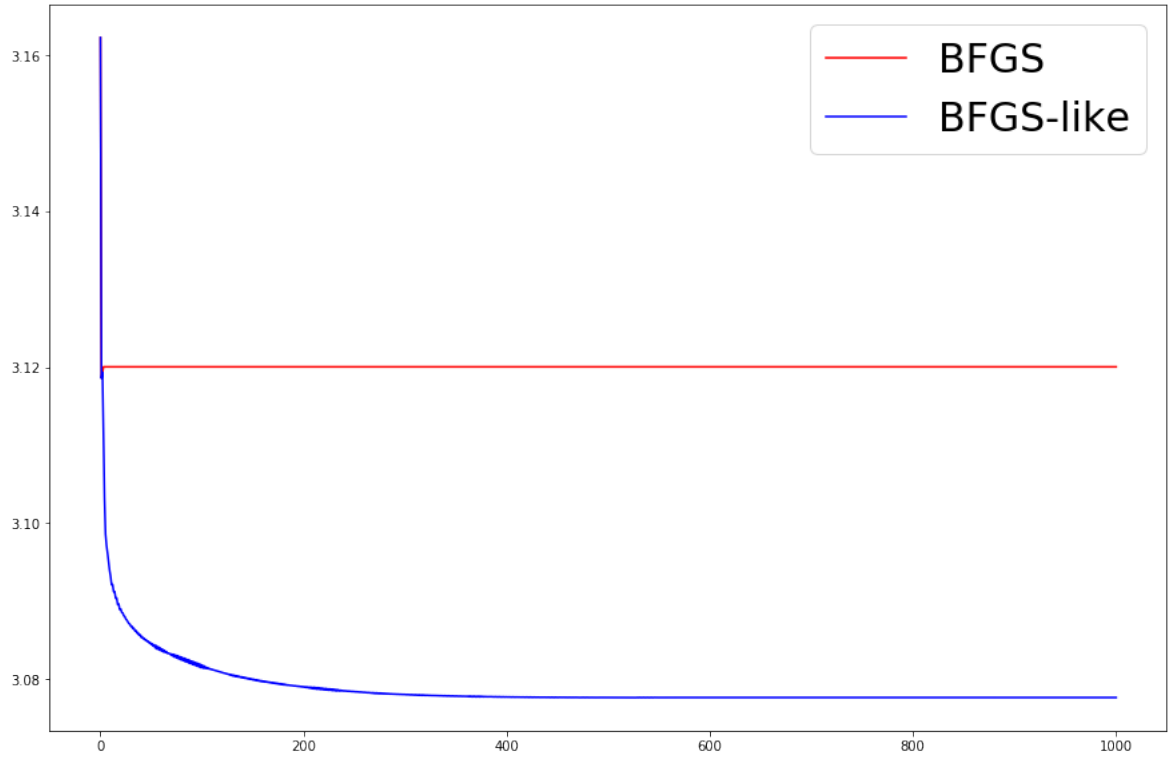
**Figure 1.** Plot of  $\|x^n - x^*\|_2$  for each iteration  $n$

Consider now the (nonconvex) Generalized Rosenbrock function in 10 dimensions:

$$f(x) = \sum_{i=1}^9 \left[ 100 (x_{i+1} - x_i^2)^2 + (x_i - 1)^2 \right], \quad -30 \leq x_i \leq 30, \quad i = 1, 2, \dots, 10.$$

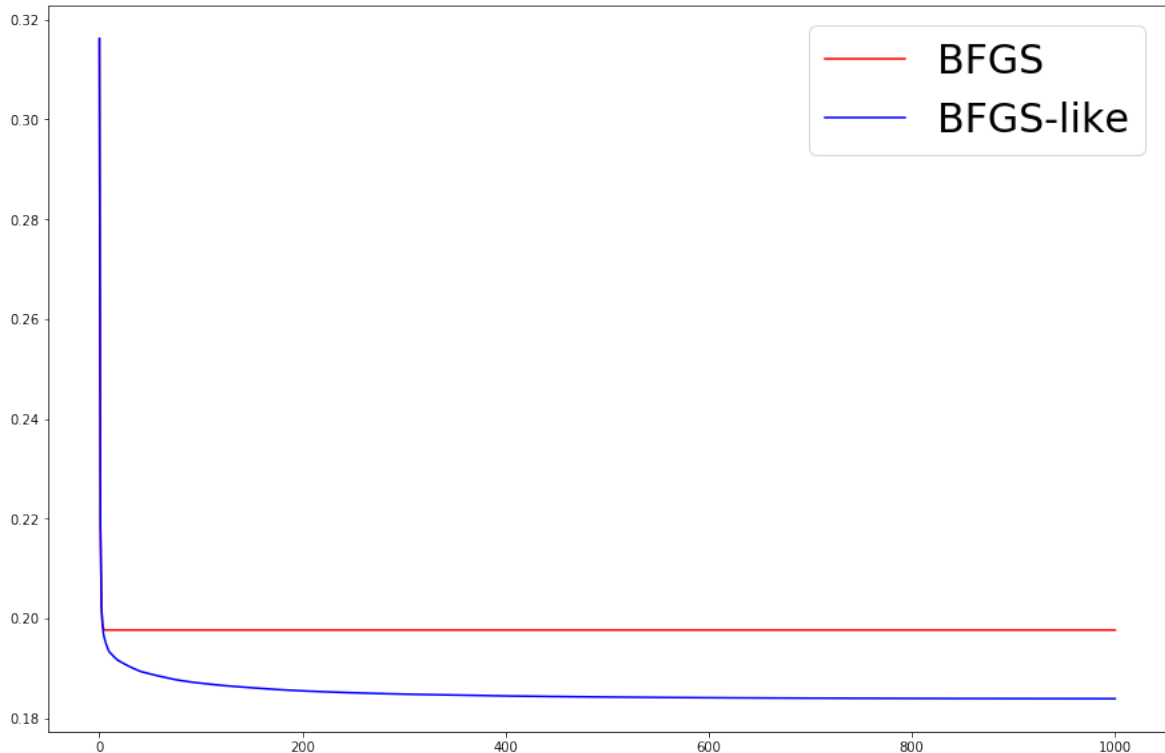
It has an absolute minimum at  $x_i^* = 1, i = 1, \dots, 10$  and  $f(x^*) = 0$ . Taking as initial point  $x_0 = (0, 0, \dots, 0)$  the origin, both methods get stuck in a local minimum, see Figure 2.





**Figure 2.** Plot of  $\|x^n - x^*\|_2$  for each iteration  $n$

Instead, initiating the recursions at  $x_0 = (0.9, 0.9, \dots, 0.9)$ , both algorithms converge to the absolute minimum (Figure 3 depicts 100 iterations). After a few initial steps, BFGS-like appears to perform better than BFGS.



**Figure 3.** Plot of  $\|x^n - x^*\|_2$  for each iteration  $n$

## 6. Closing comments

We have proposed a new family of BFGS-like iterations of which (13)-(14) is a most natural one. The entropic variational derivation provides theoretical support for these methods and a new proof of Fletcher’s classical derivation [3]. Further study is needed to exploit the flexibility afforded by this new family (the vector  $v_k$  determining the oblique projection in (15) appears as a “free parameter”). Similar results have been established for block BFGS. A few numerical experiments seem to indicate that (13)-(14) may perform better in some problems than standard BFGS.

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